# Turning Borel sets into Clopen effectively 

Vassilis Gregoriades<br>TU Darmstadt gregoriades@mathematik.tu-darmstadt.de

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Theorem. If $A$ is a Borel subset of a Polish space $(\mathcal{X}, \mathcal{T})$ there exists a Polish topology $\mathcal{T}_{\infty}$ on $\mathcal{X}$ which extends $\mathcal{T}$, and thus has the same Borel sets as $\mathcal{T}$ such that $A$ is $\mathcal{T}_{\infty}$-clopen.
Theorem. (Lusin-Suslin) Every Borel subset of a Polish space is the continuous injective image of a closed subset of the Baire space $\mathcal{N}={ }^{\omega} \omega$.
We consider the family of all recursive functions from $\omega^{k}$ to $\omega^{n}$. A set $P \subseteq \omega^{k}$ is recursive when the characteristic function $\chi_{p}$ is recursive.
Relativization. For every $\varepsilon \in \mathcal{N}$ one defines the relativized family of $\varepsilon$-recursive functions. Similarly one defines the family of $\varepsilon$-recursive subsets of $\omega^{k}$.

Definition. (Moschovakis) Suppose that $\mathcal{X}$ is a Polish space, $d$ is compatible distance function for $\mathcal{X}$ and $\left(x_{n}\right)_{n \in \omega}$ is a sequence in $\mathcal{X}$. Define the relation $P_{<}$of $\omega^{4}$ as follows
$P_{<}(i, j, k, m) \Longleftrightarrow d\left(x_{i}, x_{j}\right)<\frac{k}{m+1}$. Similarly we define the relation $P_{\leq}$.
The sequence $\left(x_{n}\right)_{n \in \omega}$ is a recursive presentation of $\mathcal{X}$, if
(1) it is a dense sequence and
(2) the relations $P_{<}$and $P_{\leq}$are recursive.

The spaces $\mathbb{R}, \mathcal{N}$ and $\omega^{k}$ admit a recursive presentation i.e., they are recursively presented. Some other examples: $\mathbb{R} \times \omega, \mathbb{R} \times \mathcal{N}$. However not all Polish spaces are recursively presented.
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Every Polish space admits an $\varepsilon$-recursive presentation for some suitable $\varepsilon$.
$N(\mathcal{X}, s)=$ the ball with center $x_{(s)_{0}}$ and radius $\frac{(s)_{1}}{(s)_{2}+1}$.
A set $P \subseteq \mathcal{X}$ is semirecursive if $P=\bigcup_{i \in \omega} N(\mathcal{X}, \alpha(i))$ where $\alpha$ is a recursive function from $\omega$ to $\omega$.
$\Sigma_{1}^{0}=$ all semirecursive sets
$\rightsquigarrow$ effective open sets.
$\Pi_{1}^{0}=$ the complements of semirecursive sets
$\rightsquigarrow$ effective closed sets.
Similarly one defines the class $\Delta_{1}^{1}$ of effective Borel sets, $\Sigma_{1}^{1}$ of effective analytic and so on.
A function $f: \mathcal{X} \rightarrow \mathcal{Y}$ is $\Sigma_{1}^{0}$-recursive if and only if the set $R^{f} \subseteq \mathcal{X} \times \omega, R^{f}(x, s) \Longleftrightarrow f(x) \in N(\mathcal{Y}, s)$, is $\Sigma_{1}^{0}$.
A point $x \in \mathcal{X}$ is $\Delta_{1}^{1}$ point if the relation $U \subseteq \omega$ which is defined by $s \in U \Longleftrightarrow x \in N(\mathcal{X}, s)$, is in $\Delta_{1}^{1}$.
Similarly one defines the relativized pointclasses with respect to some parameter $\varepsilon$.
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Similarly one defines the relativized pointclasses with respect to some parameter $\varepsilon$.

Theorem. Every $\Delta_{1}^{1}$ subset of a recursively presented Polish space is the recursive injective image of a $\Pi_{1}^{0}$ subset of $\mathcal{N}$.
Theorem. (G.) Suppose that $(\mathcal{X}, \mathcal{T})$ is a recursively presented Polish space, $d$ is a suitable distance function for $(\mathcal{X}, \mathcal{T})$ and $A$ is a $\Delta_{1}^{1}$ subset of $\mathcal{X}$. There exists an $\varepsilon_{A} \in \mathcal{N}$, which is recursive in Kleene's $O$ and a Polish topology $\mathcal{T}_{\infty}$ with suitable distance function $d_{\infty}$, which extends $\mathcal{T}$ and has the following properties:
(1) The Polish space $\left(\mathcal{X}, \mathcal{T}_{\infty}\right)$ is $\varepsilon_{A}$-recursively presented.
(2) The set $A$ is a $\Delta_{1}^{0}\left(\varepsilon_{A}\right)$ subset of $\left(\mathcal{X}, d_{\infty}\right)$.
(3) If $B \subseteq \mathcal{X}$ is a $\Delta_{1}^{1}(\alpha)$ subset of $(\mathcal{X}, d)$, where $\alpha \in \mathcal{N}$, then $B$ is a $\Delta_{1}^{1}\left(\varepsilon_{A}, \alpha\right)$ subset of $\left(\mathcal{X}, d_{\infty}\right)$.
(4) If $B \subseteq \mathcal{X}$ is a $\Delta_{1}^{1}\left(\varepsilon_{A}, \alpha\right)$ subset of $\left(\mathcal{X}, d_{\infty}\right)$, where $\alpha \in \mathcal{N}$, then $B$ is a $\Delta_{1}^{1}\left(\varepsilon_{A}, \alpha\right)$ subset of $(\mathcal{X}, d)$.

Remark. If the inverse function in the Lusin-Suslin Theorem is continuous, then the set $A$ that we start with is $G_{\delta}$.
Lemma. (G.) For every $A \subseteq \mathcal{N}$ in $\Pi_{2}^{0}$ there is a set $F \subseteq \mathcal{N}$ in $\Pi_{1}^{0}$ and a recursive function $\pi: \mathcal{N} \rightarrow \mathcal{N}$ which is injective on $A$ such that $\pi[F]=A$ and the inverse $\pi^{-1}$ is continuous.
Corollary. (G.) Suppose that $A$ is a $\Delta_{1}^{1}$ subset of $\mathcal{N}$, which is also in ${\underset{\sim}{~}}_{2}^{0}$ and assume moreover that the class $\Delta_{1}^{1}$ is dense in $A$ and $\mathcal{N} \backslash A$. Then one can choose the previous parameter $\varepsilon_{A}$ in $\Delta_{1}^{1}$. Sketch of the proof. It's just a sketch - really! From of a theorem of Louveau the set $A$ is in $\Delta_{2}^{0}(\varepsilon)$ for some $\varepsilon \in \Delta_{1}^{1}$. Apply the previous lemma and proceed as usual.
Theorem (The Strong $\Delta$-Selection Principal). Suppose that $\mathcal{Z}$ and $\mathcal{Y}$ are recursively presented Polish spaces and that $P \subseteq \mathcal{Z} \times \mathcal{Y}$ is in $\Pi_{1}^{1}$ and such that for all $z \in \mathcal{Z}$ there exists $y \in \Delta_{1}^{1}(z)$ such that $(z, y) \in P$. Then there exists a $\Delta_{1}^{1}$-recursive function $f: \mathcal{Z} \rightarrow \mathcal{Y}$ such that $(z, f(z)) \in P$ for all $z \in \mathcal{Z}$.

Corollary. (G.) Suppose that $\mathcal{Z}$ is a Polish space, $\mathcal{X}$ is a closed subset of $\mathcal{N}$ and that $P$ is a Borel subset of $\mathcal{Z} \times \mathcal{X}$ such that the sets $P_{z}$ and $\mathcal{X} \backslash P_{z}$ are infinite for all $z \in \mathcal{Z}$. Assume moreover that $(*) \Delta_{1}^{1}(z)$ is dense in both $P_{z}$ and $\mathcal{X} \backslash P_{z}$ for all $z \in \mathcal{Z}$.
Then there is a Borel-measurable function $f: \mathcal{Z} \rightarrow \mathcal{N}$ such that $f(z)$ "encodes" a distance function $d_{z}$ on $\mathcal{X}$ such that: (1) the space $\left(\mathcal{X}, d_{z}\right)$ is complete and separable, (2) the topology $\mathcal{T}_{d_{z}}$ extends $\mathcal{T}$ and (3) $P_{z}$ is $d_{z}$-clopen, for all $z \in \mathcal{Z}$.
Thanks to results of Tanaka, Sacks, Thomason and Hinman, we may replace the effective condition $(*)$ with one of the following classical conditions:
(1) there is a "reasonable" Borel measure $\mu$ on $\mathcal{X}$ such that for all open $V$ and for all $z \in \mathcal{Z}$ if $P_{z} \cap V \neq \emptyset$ we have that $P_{z} \cap V$ is countable or $\mu\left(P_{z} \cap V\right)>0$. Similarly for $\mathcal{X} \backslash P_{z}$;
(2) $P_{z}$ is countable or co-countable for all $z \in \mathcal{Z}$.

